

Appendix B

ALGORITHMS

Here we adopt a standard notation (pseudocode) for defining algorithms [33, 154]. Assignment of a value α to a variable β is denoted by $\beta \leftarrow \alpha$. The branching probability $P_{\alpha,\beta}$ is represented as $P[\alpha, \beta]$ if it is stored in matrix form, and as $P[\beta][\alpha]$ if it is stored in the form of adjacency lists.

Algorithm B.1 Calculate the pathway sum $\mathcal{S}_{\alpha,\beta}^{C_N}$

Require: Chain nodes are numbered $0, 1, 2, \dots, N - 1$
Require: $1 < N$
Require: $\alpha, \beta \in \{0, 1, 2, \dots, N - 1\}$
Require: $P[i, j]$ is the probability of branching from node j to node i

- 1: **if** $\alpha < \beta$ **then**
- 2: $h \leftarrow 1; t \leftarrow N - 2; s \leftarrow 1$
- 3: **else**
- 4: $h \leftarrow N - 2; t \leftarrow 1; s \leftarrow -1$
- 5: **end if**
- 6: $L \leftarrow 1$
- 7: **for all** $i \in \{h, h + s, h + 2s, \dots, \alpha\}$ **do**
- 8: $L \leftarrow 1/(1 - P[i - s, i]P[i, i - s]L)$
- 9: **end for**
- 10: $\Pi \leftarrow 1$
- 11: **for all** $i \in \{\alpha + s, \alpha + 2s, \alpha + 3s, \dots, \beta\}$ **do**
- 12: $\Pi \leftarrow P[i - s, i]L\Pi$
- 13: $L \leftarrow 1/(1 - P[i - s, i]P[i, i - s]L)$
- 14: **end for**
- 15: $R \leftarrow 1$
- 16: **for all** $i \in \{t, t - s, t - 2s, \dots, \beta\}$ **do**
- 17: $R \leftarrow 1/(1 - P[i + s, i]P[i, i + s]R)$
- 18: **end for**
- 19: **return** $LR\Pi/(L - LR + R)$

Algorithm B.2 Calculate the pathway sum $\mathcal{S}_{a,b}^{G_N}$

Require: $1 < N$
Require: $a, b \in \{0, 1, 2, \dots, N - 1\}$
Require: W is a boolean array of size N with every element initially set to True
Require: N_W is the number of True elements in array W (initialised to N)
Require: $P[i, j]$ is the probability of branching from node j to node i
Require: $AdjIn[i]$ and $AdjOut[i]$ are the lists of indices of all nodes connected to node i via incoming and outgoing edges, respectively

Recursive function $F(\alpha, \beta, W, N_W)$

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1:  $W[\beta] \leftarrow \text{False}$ 
2:  $N_W \leftarrow N_W - 1$ 
3: if  $\alpha = \beta$  and  $N_W = 0$  then
4:    $\Sigma \leftarrow 1$ 
5: else
6:    $\Sigma \leftarrow 0.0$ 
7:   for all  $i \in AdjOut[\beta]$  do
8:     for all  $j \in AdjIn[i]$  do
9:       if  $W[i]$  and  $W[j]$  then
10:         $\Sigma \leftarrow \Sigma + P[\beta, j]F(j, i, W, N_W)P[i, \beta]$ 
11:       end if
12:     end for
13:   end for
14:    $\Sigma \leftarrow 1/(1 - \Sigma)$ 
15:   if  $\alpha \neq \beta$  then
16:      $\Lambda \leftarrow 0.0$ 
17:     for all  $i \in AdjOut[\beta]$  do
18:       if  $W[i]$  then
19:          $\Lambda \leftarrow \Lambda + F(\alpha, i, W, N_W)P(i, \beta)$ 
20:       end if
21:     end for
22:      $\Sigma \leftarrow \Sigma \Lambda$ 
23:   end if
24: end if
25:  $W[\beta] \leftarrow \text{True}$ 
26:  $N_W \leftarrow N_W + 1$ 
27: return  $\Sigma$ 

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Algorithm B.3 Calculate the pathway sum $S_{\alpha,\beta}^{\mathcal{G}_N}$ from every source to every sink, and the mean escape time for every source in a dense graph \mathcal{G}_N

Require: Nodes are numbered $0, 1, 2, \dots, N - 1$
Require: Sink nodes are indexed first, source nodes - last
Require: i is the index of the first intermediate node
Require: s is the index of the first source node
Require: In case there are no intermediate nodes $i = s$, otherwise $i < s$
Require: $1 < N$
Require: $i, s \in \{0, 1, 2, \dots, N - 1\}$
Require: $\tau[\alpha]$ is the waiting time for node α , $\alpha \in \{i, i + 1, i + 2, \dots, N - 1\}$
Require: $P[i, j]$ is the probability of branching from node j to node i

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1: for all  $\gamma \in \{i, i + 1, i + 2, \dots, s - 1\}$  do
2:   for all  $\beta \in \{\gamma + 1, \gamma + 2, \dots, N - 1\}$  do
3:     if  $P[\gamma, \beta] > 0$  then
4:        $\tau[\beta] \leftarrow (\tau[\beta] + \tau[\gamma]P[\gamma, \beta])/(1 - P[\beta, \gamma]P[\gamma, \beta])$ 
5:       for all  $\alpha \in \{0, 1, 2, \dots, N - 1\}$  do
6:         if  $\alpha \neq \beta$  and  $\alpha \neq \gamma$  then
7:            $P[\alpha, \beta] \leftarrow (P[\alpha, \beta] + P[\alpha, \gamma]P[\gamma, \beta])/(1 - P[\beta, \gamma]P[\gamma, \beta])$ 
8:         end if
9:       end for
10:       $P[\gamma, \beta] \leftarrow 0.0$ 
11:    end if
12:  end for
13:  for all  $\alpha \in \{0, 1, 2, \dots, N - 1\}$  do
14:     $P[\alpha, \gamma] \leftarrow 0.0$ 
15:  end for
16: end for
17: for all  $\alpha \in \{s, s + 1, s + 2, \dots, N - 1\}$  do
18:   for all  $\beta \in \{s, s + 1, s + 2, \dots, N - 1\}$  do
19:     if  $\alpha \neq \beta$  and  $P[\alpha, \beta] > 0$  then
20:        $P_{\alpha,\beta} \leftarrow P[\alpha, \beta]$ 
21:        $P_{\beta,\alpha} \leftarrow P[\beta, \alpha]$ 
22:        $T \leftarrow \tau[\alpha]$ 
23:        $\tau[\alpha] \leftarrow (\tau[\alpha] + \tau[\beta]P_{\beta,\alpha})/(1 - P_{\alpha,\beta}P_{\beta,\alpha})$ 
24:        $\tau[\beta] \leftarrow (\tau[\beta] + TP_{\alpha,\beta})/(1 - P_{\alpha,\beta}P_{\beta,\alpha})$ 
25:       for all  $\gamma \in \{0, 1, 2, \dots, i - 1\} \cup \{s, s + 1, s + 2, \dots, N - 1\}$  do
26:          $T \leftarrow P[\gamma, \alpha]$ 
27:          $P[\gamma, \alpha] \leftarrow (P[\gamma, \alpha] + P[\gamma, \beta]P_{\beta,\alpha})/(1 - P_{\alpha,\beta}P_{\beta,\alpha})$ 
28:          $P[\gamma, \beta] \leftarrow (P[\gamma, \beta] + TP_{\alpha,\beta})/(1 - P_{\alpha,\beta}P_{\beta,\alpha})$ 
29:       end for
30:        $P[\alpha, \beta] \leftarrow 0.0$ 
31:        $P[\beta, \alpha] \leftarrow 0.0$ 
32:     end if
33:   end for
34: end for

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Algorithm B.4 Detach node γ from an arbitrary graph \mathcal{G}_N

Require: $1 < N$
Require: $\gamma \in \{0, 1, 2, \dots, N - 1\}$
Require: $\tau[i]$ is the waiting time for node i
Require: $Adj[i]$ is the ordered list of indices of all nodes connected to node i via outgoing edges
Require: $|Adj[i]|$ is the cardinality of $Adj[i]$
Require: $Adj[i][j]$ is the index of the j th neighbour of node i
Require: $P[i]$ is the ordered list of probabilities of leaving node i via outgoing edges, $|P[i]| = |Adj[i]|$
Require: $P[i][j]$ is the probability of branching from node i to node $Adj[i][j]$

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1: for all  $\beta_\gamma \in \{0, 1, 2, \dots, |Adj[\gamma]| - 1\}$  do
2:    $\beta \leftarrow Adj[\gamma][\beta_\gamma]$ 
3:    $\gamma_\beta \leftarrow -1$ 
4:   for all  $i \in \{0, 1, 2, \dots, |Adj[\beta]| - 1\}$  do
5:     if  $Adj[\beta][i] = \gamma$  then
6:        $\gamma_\beta \leftarrow i$ 
7:       break
8:     end if
9:   end for
10:  if not  $\gamma_\beta = -1$  then
11:     $P_{\beta,\beta} \leftarrow 1/(1 - P[\beta][\gamma_\beta]P[\gamma][\beta_\gamma])$ 
12:     $P_{\beta,\gamma} \leftarrow P[\beta][\gamma_\beta]$ 
13:     $Adj[\beta] \leftarrow \{Adj[\beta][0], Adj[\beta][1], \dots, Adj[\beta][\gamma_\beta - 1], Adj[\beta][\gamma_\beta + 1], \dots\}$ 
14:     $P[\beta] \leftarrow \{P[\beta][0], P[\beta][1], \dots, P[\beta][\gamma_\beta - 1], P[\beta][\gamma_\beta + 1], \dots\}$ 
15:    for all  $\alpha_\gamma \in \{0, 1, 2, \dots, |Adj[\gamma]| - 1\}$  do
16:       $\alpha \leftarrow Adj[\gamma][\alpha_\gamma]$ 
17:      if not  $\alpha = \beta$  then
18:        if exists edge  $\beta \rightarrow \alpha$  then
19:          for all  $i \in \{0, 1, 2, \dots, |Adj[\beta]| - 1\}$  do
20:            if  $Adj[\beta][i] = \alpha$  then
21:               $P[\beta][i] \leftarrow P[\beta][i] + P_{\beta,\gamma}P[\gamma][\alpha_\gamma]$ 
22:              break
23:            end if
24:          end for
25:        else
26:           $Adj[\beta] \leftarrow \{\alpha, Adj[\beta][0], Adj[\beta][1], Adj[\beta][2], \dots\}$ 
27:           $P[\beta] \leftarrow \{P_{\beta,\gamma}P[\gamma][\alpha_\gamma], P[\beta][0], P[\beta][1], P[\beta][2], \dots\}$ 
28:        end if
29:      end if
30:    end for
31:    for all  $i \in \{0, 1, 2, \dots, |P[\beta]| - 1\}$  do
32:       $P[\beta][i] \leftarrow P[\beta][i]P_{\beta,\beta}$ 
33:    end for
34:     $\tau_\beta \leftarrow (\tau_\beta + P_{\beta,\gamma}\tau_\gamma) P_{\beta,\beta}$ 
35:  end if
36: end for

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